# Randomness in Scientific Computing 

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## Today's class

Today we will discuss:

- Randomness, why you want it.
- How to make it or fake it.
- Applications: Monte Carlo


## Why Randomness?

## Why Randomness?

- To simulate some physical phenomenon that has noise.
E.g. Brownian motion, Nyquist noise.

On the level of their description, this is real randomness.

- To perform averages or integrals in systems with many degrees of freedom.
E.g. Stat. Phys. computations, path integral calculations.

Here, the main objective is to get the converged answer quickly.

- To estimate a parameter's distribution from using data (MCMC).
- To test a statistical method.


## Creating Randomness

## Sources of randomness

## True Random Number Generators

- Lava lamps.
- Radioactive decay.
- Various quantum processes.
- Atmospheric noise.
- Random computer hardware noise signals (thermals noise).

Generally slow, expensive, impossible to reproduce for debugging. Hard to characterize underlying distribution.

## Pseudo Random Number Generators

- Come up with a algorithm that produces random numbers
- But wouldn't such an algorithm would be deterministic?
- Only has to act random, i.e., give fair and uncorrelated sequence.


## Pseudo Random Number Generators (PRNG)

Recipe:

- Define some 'state', initialized by some 'seed' value(s).
- Produce a number from this state.
- Advance the state determistically.
- As long as the numbers produces behave as if they are
- independent
> identically distributed
- according to a predefined distribution (eg uniform)
we will be satisfied.
Depends a lot on the way the states are advanced. Must test.


## Distributions are transformations

- Suppose we had a way to draw random values of a continuous variable $\boldsymbol{x}$ that is uniformly distributed between 0 and 1.
- Let's say that for any value x that is drawn, we were to compute a value $y=f(x)$, where $f$ is a deterministic function.
- The values of $\boldsymbol{y}$ are also randomly distributed, but with a non-uniform distribution (unless $f(x)=x$ ).

So we can turn a uniformly distributed random variable into a non-uniformly distributed variable by applying a function.

If we want a specific non-uniform distribution, we just need to figure out the function. For many common cases, this is already done.

So our main focus is first to find uniformly distributed variables.

## All pseudo random numbers are discrete

Despite the illusion of continous variables that floating point numbers give, there are only a finite number of bits, and thus a discrete set of values.

In fact, routines that give pseudo random floating point numbers are usually based on drawing a random integer number and dividing it by the largest possible generated integer.

From a random integer of $n$ bits, we just need each bit to be uniformly distributed, with a chance of $50 \%$ of a 0 and $50 \%$ of a 1 .

Warning: most PRNGs give lower bits that are more correlated than the higher bits.

## Example: Coin Toss

The following class can produce a 'random' 1's and 0's representing heads and tails:

```
// badcoin.h
class BadCoin {
    public:
        // method to set the starting seed
    void start(unsigned int seed) {
        state = seed;
    }
    // method to toss the coin (1:head, 0:tail)
    int toss() {
        state++; // update state
        return state%2; // using lowest bit...
    }
    private:
    unsigned int state; // internal state
};
```

- Is it fair? Independent samples? Period?

```
#include <iostream>
#include "badcoin.h"
int main()
{
    BadCoin coin;
    coin.start(13); //seed
    // toss the coin 20 times
    for (int i = 0; i < 20; i++)
        std::cout << coin.toss() << '\n';
    return 0;
}
```

What does this give?

## Testing for randomness

Suppose we have drawn $N$ samples using our PRNG.
Let's look at two tests:
(1) Fairness: histogram counting the occurance of values

$$
h_{x}=\sum_{i=1}^{N} \delta_{x x_{i}}
$$

Here $x$ is one of the possible random numbers (here $\pm 1$ ), and $x_{i}$ are samples produced by our PRNG $\left(\delta_{i i}=1, \delta_{i, j \neq i}=0\right)$.
(2) Independence: look at correlations between samples:

$$
c_{j}=\left\langle x_{i} x_{i+j}\right\rangle=\frac{1}{N} \sum_{i=1}^{N}\left(x_{i}-\bar{x}\right)\left(x_{i+j}-\bar{x}\right)
$$

If independent: $\mathcal{O}(1 / \sqrt{N})$ for $j>0$

## Test results $(\mathrm{N}=20)$

Independence


Bad!

## Try again

```
Old version
// badcoin.h
class BadCoin {
    public:
    // method to set the starting seed
    void start(int seed) {
        state = seed;
    }
    // method to toss the coin (1:head, 0:tail)
    int toss() {
        state++; // update state
        return state%2; // using lowest bit...
    }
private:
    unsigned int state; // internal state
};
```


## New version

```
```

// improvedcoin.h

```
```

// improvedcoin.h
class ImprovedCoin {
class ImprovedCoin {
public:
public:
// method to set the starting seed
// method to set the starting seed
void start(int seed) {
void start(int seed) {
state = seed;
state = seed;
}
}
// method to toss the coin (1:head, 0:tail)
// method to toss the coin (1:head, 0:tail)
int toss() {
int toss() {
state=100+100*sin(state+1);//update state
state=100+100*sin(state+1);//update state
return state%2; // using lowest bit...
return state%2; // using lowest bit...
}
}
private:
private:
unsigned int state;
unsigned int state;
};

```
```

};

```
```

Difference lies in a more complex update of the state.

## "Improved" test results $(\mathbf{N}=20)$




Better?

## Let's do more samples: $N=200$

Independence


## Common RNG Types

## Linear Congruential generators

$$
x_{i+1}=\left(a x_{i}+c\right) \bmod m
$$

The quality of the random numbers depends on the parameters ( $a, c, m$ ).
Even the best ones are not very good, but they can be used as part of better generators.

## Lagged-Fibonacci generator

$$
x_{i}=\left(x_{i-j} \circ x_{i-k}\right) \bmod m
$$

where o can be any binary operator (add, mult, ...). Requires a seed block from another PRNG.

## Mersenne Twister

A complex variation of lagged-Fibonacci that stikes a great balance between speed and statistical tests.

## Other tests

- Check moments of distributions.
- Check that spacings between random points follow a Poisson integral if uniformly distributed.
- Examine sequences of 5 numbers. There are 120 ways to sort 5 numbers. The 120 ways should occur with equal probability.
- Parking circle test: randomly place unit circles in a $100 \times 100$ square. If the circle overlaps an existing one, try again. After 12,000 tries, the number of successfully "parked" ' circles should follow a certain normal distribution.
- Play 200,000 games of a dice game, counting the wins and number of throws per game. Each count follow a certain distribution.
- And many others. See, for example, the NIST test suite: https://csrc.nist.gov/projects/random-bit-generation and the TestU01 suite:
http://simul.iro.umontreal.ca/testu01/tu01.html


## Lesson: Don't do it yourself

What properties do we expect from a random number generator?

- We would like them from a given distribution (uniform, Gaussian).
- We would like them to be unpredictable.
- We would like them to be reproducible.
- We need them to be generated quickly.
- We need to have a long period.

It is not that easy to guess good PRNG algorithms and parameters.
There was a time when one was forced to implement PRNGs oneself, as standard ones were quite bad, but $\mathrm{C}++$ has decent random number generators in its <random> standard library.

## Using existing random numbers

```
Previous way
// improvedcoin.h
class ImprovedCoin {
    public:
    // method to set the starting seed
    void start(int seed) {
        state = seed;
    }
    // method to toss the coin (1:head, 0:tail)
    int toss() {
        state = 100+100*sin(state+1); // update state
        return state%2; // using lowest bit...
    }
private:
    unsigned int state;
};
```

```
C++ way
// goodcoin.h
#include <random>
class GoodCoin {
    public:
        GoodCoin(): uniform(0,1) {}
    // method to set the starting seed
    void start(int seed) {
        engine.seed(seed);
    }
    // method to toss the coin (1:head, 0:tail)
    int toss() {
        return uniform(engine);//state in engine
    }
private:
    std::uniform_int_distribution<int> uniform;
    std::mt19937 engine; // PNRG state
};
```


## Test C++ way, N=200



Independence

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## About the random standard libary

The <random> library allows to produce random numbers using combinations of generators and distributions.

Generators Objects that generate uniformly distributed numbers.
Distributions Objects that transform sequences of numbers generated by a generator into sequences of numbers that follow a specific random variable distribution, such as uniform, Normal or Binomial.

Distribution objects generate random numbers by means of their operator() member, which takes a generator object as argument:

```
std::mt19937_64 generator;
std::uniform_int_distribution<int> distribution(1,6);
int die_roll = distribution(generator); // generates number in the range 1..6
```


## Available generators

While there are ways to create your own, the library has a number of standard available generators:

| default_random_engine | Default random engine |
| :--- | :---: |
| minstd_rand | Minimal Standard minstd_rand generator |
| minstd_rand0 | Minimal Standard minstd_rand0 generator |
| mt19937 | Mersenne Twister 19937 generator |
| mt19937_64 | Mersenne Twister 19937 generator (64 bit) |
| ranlux24_base | Ranlux 24 base generator |
| ranlux48_base | Ranlux 48 base generator |
| ranlux24 | Ranlux 24 generator |
| ranlux48 | Ranlux 48 generator |
| knuth_b | Knuth-B |

## Some good PRNGs

The following have long periods, independent samples, a fair distribution, and pass most statistical tests:

- Mersenne twister: mt19937 and mt19937_64, in the C++ random library. Use this one if you need many billions of random numbers relatively fast.
- In the lagged-Fibonacci class: ranlux24 and ranlux48, in the C++ random library. Use this if speed is not an impediment and you need more statistical tests passed.
- r1279 (lagged-Fibonacci generator), in the GSL and boost::random.
- WELL generator
https://www.arxiv-vanity.com/papers/1005.4117


## Tip

Employ two random number generators, and see if they give, statistically speaking, the same result. If they don't, one of them is bad for your application.

## Monte Carlo

## Monte Carlo Techniques

A collection of techniques whose unifying feature is the use of randomness. These applications of randomness generally fall into one of three categories:

- Adding randomness to otherwise-deterministic dynamics, and studying how the dynamics are changed.
- Generating samples from a given probability distribution, $\boldsymbol{P}(\boldsymbol{x})$, usually a distribution that is complicated and can't be dealt with nicely in closed form (e.g. Markov Chain Monte Carlo).
- Estimating expectation values under this distribution, e.g.

$$
\langle A(\mathrm{x})\rangle=\int P(\mathrm{x}) A(\mathrm{x}) d \mathrm{x}
$$

where $\mathbf{x}$ is typically high dimensional.

```
These depend on having a good random number generator!
```


## Monte Carlo example: traffic flow

Nagel-Schreckenberg traffic is a 1D toy model used to generate traffic-like behaviour. At each time step in the model, the following rules are applied to each car in the simulation:
(1) If the velocity is below vmax, then increase $v$ by 1 (try to speed up).

2 If the car in front of the given car is a distance $d$ away, and $v \geq d$, then reduce $v$ to $d-1$ (don't want to hit the car).
(3) Add randomness: if $\mathrm{v}>0$ then with probability p the car reduces its speed by 1 .
(4) The car moves ahead by v steps (on a circular track).

The four rules boil down to

$$
\begin{aligned}
& v \leftarrow \min \left(v+1, v_{\max }\right) \\
& v \leftarrow \min (v, d-1) \\
& v \leftarrow v-1 \text { if } v \neq 0 \text { with probability } p \\
& x \leftarrow x+v
\end{aligned}
$$

## Monte Carlo example: traffic flow

numcars $=200$
gridsize $=1000$
$\mathrm{p}=0.13$
$\operatorname{vmax}=5$


## Note on Implementing Chance

$$
v \leftarrow v-1 \text { if } v \neq 0 \text { with probability } p
$$

How do you do that?

- Draw a random number $r$ using a PRNG with uniform distribution on $[0,1)$.
- For any chosen value $p \in[0,1)$, the chance that $r$ is less than that value, is $p$ itself.
- So if $\boldsymbol{r}$ is less than $\boldsymbol{p}$, we will accept the move and decrease $\boldsymbol{v}$ if possible.
- If $\boldsymbol{r}$ is greater than or equal to $\boldsymbol{p}$, we leave $\boldsymbol{v}$ as it is, i.e., we reject the move.


## Monte Carlo Example: Chemical bond breakage



## Model parameters

- the initial bond length
- the temperature (sets strength of thermal fluctuations)
- In this model, the molecular bond between the atoms can break due to thermal fluctuations.
- Escape from the potential well corresponds to the bond breaking.


## Simulation parameters

- the timestep
- maximum time to simulate
- random seed
- plus output parameters


## Monte Carlo Example: Dynamics implementation

For each "timestep":
(1) Randomly perturb the bond extension, proportional to $\Delta t$.

2 Calculate the new energy of the system.

- If the energy of the system goes down, keep the new position.
- If the energy of the system goes up, keep the position if $r<\exp (-\Delta E / T)$, where $r$ is a random number between 0 and 1 , and T is the system temperature.

3 Repeat for all timesteps until escape.

## Monte Carlo Example: results


initial bond extension: 0.6
temperature: 2.8
timestep: 0.0003
seed: 13

Breakage occured at $\mathrm{t}=5.334000000001951$.

